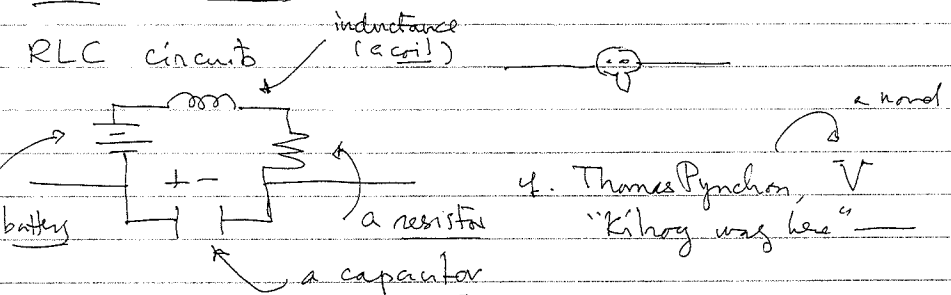


Starting or initial transients:

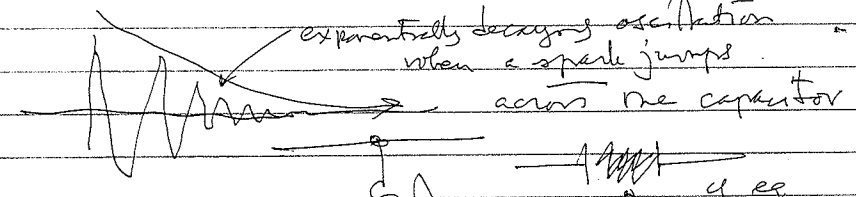


$I = I(t) = \text{current (in amperes)}$

through the circuit at time  $t$   
satisfies a 2<sup>nd</sup> order differential equation

$L \cdot \frac{d^2 I}{dt^2} + R \cdot \frac{dI}{dt} + C \cdot I = 0$	$L = \text{inductance}$ $C = \text{capacitance}$ $R = \text{resistance}$ $\geq 0$ <u>constants</u>
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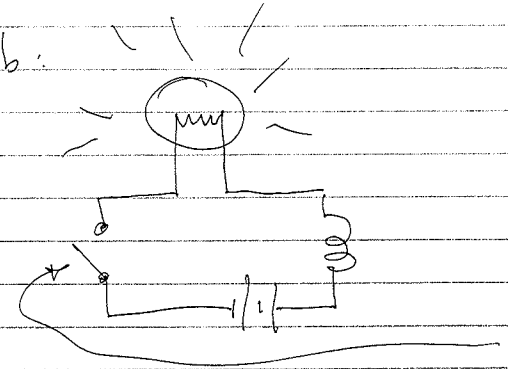
with a solution  $I(t)$



Formally the same eqn as the hanging spring  
 e.g. thunder & lightning

ii)

This is pretty complicated, so let's simplify to the case of an old-fashioned incandescent lightbulb:



No capacitor, but instead we have a switch

which we turn on at  $t=0$ : thus we

have the system of equations

$$L \frac{d^2 I}{dt^2} + R \frac{dI}{dt} + \cancel{C \cdot I} = 0$$

and initial conditions  $I(t) = 0$  at  $t=0$

To simplify, let  $d = \frac{dI}{dt}$ : then we have

$$L \frac{dd}{dt} + R d = 0 \Leftrightarrow L dd + R d dt = 0$$

iii)

Divide by  $I$ , and integrate:

$$\frac{dI}{I} = -\frac{R}{L} dt \Rightarrow$$

$\uparrow$   
 constant

$$\int \frac{dI}{I} = \ln I = -\frac{R}{L} t + \text{Constant of integration}$$

Exponentiate to get out of the logarithm  
 $\Rightarrow$

$$\frac{dI}{I} = -\frac{R}{L} dt = \exp\left(-\frac{R}{L} t + \text{Const}\right)$$

$$= A e^{-R/L t} \quad \text{for some } A = e^{\text{const}}$$

Integrate again to get

$$I(t) = \int A e^{-R/L t} dt = -\frac{AL}{R} e^{-R/L t} + \text{Const.}$$

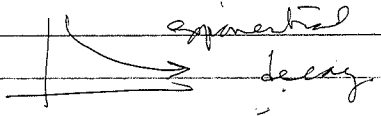
$\downarrow$   
 constant of integration

$$= B(1 - e^{-R/L t}) \quad \text{where } B = \frac{AL}{R}$$

$\uparrow$   
 rearrange

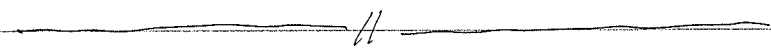
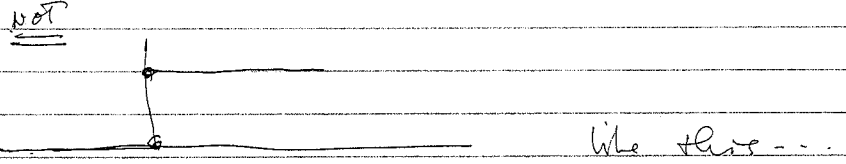
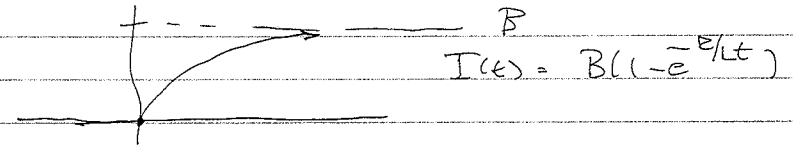
$\downarrow$   
 some constant

$\rightarrow B$  as  $t \rightarrow \infty$ , because  $e^{-R/L t} \rightarrow 0$  as  $t \rightarrow \infty$



iv)

The MORAL of the story is that the lightbulb doesn't light up instantaneously, or discontinuously:



v)

Limits

Ch 8 133-135  
Ch 29 588-600

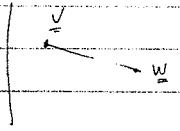
Primitives  $\mathbb{R}, \mathbb{R}^2, \mathbb{R}^3$

line, plane, 3D

sufficient for geometry, physics, real world:

but: where does  $\mathbb{R}$  come from?

In particular:



$$\text{dist}(v, w) = |v - w|$$
$$= \sqrt{(v_1 - w_1)^2 + (v_2 - w_2)^2}$$

How do we know square roots exist?

Revisiting decimal #'s:

finite repres

1.4 14 ...

$\sqrt{2} \ 00 \ 00 \ 00$

$\begin{array}{r} 4 \\ 20 \end{array} \overline{) 100}$

$\begin{array}{r} 280 \\ 2820 \end{array} \overline{) 11900}$

$\begin{array}{r} 11296 \\ 60400 \end{array}$

22 #14.e  
p. 464

$$d = d_0 \cdot d_1 d_2 \dots$$
$$d_i \in \{0, \dots, 9\}$$
$$d_0 + \frac{d_1}{10} + \frac{d_2}{10^2} + \dots + \frac{d_n}{10^n}$$

$\in$  rational numbers:  $\mathbb{Q}$

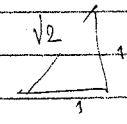
$$\frac{10^n d_0 + 10^{n-1} d_1 + \dots + d_n}{10^n}$$

= ratio of integers.

Claim: a real number is the limit of (for increasing) sequence of rational numbers.

ii)

Claim



Rational #'s are inadequate for geometry.

There is no rational number whose square is 2.

[Big problem for the Greeks]

Proof: if found then of arithmetic:

any integer  $\geq 1$  is the product of a unique set (up to order) of prime integers  $> 1$ .

$n$  is prime, or not. If it's prime, we're finished. If it's not prime, then it's divisible by something  $< n$ .  
 $n = p \cdot q$  with  $p, q < n$ .

Induction: Suppose we've proved this for all #'s smaller than  $n$ . Then  $p, q$  are products of primes, etc.

√(n)

√(n)

Claim no prime # can have a rational square root

$$p = r^2, \quad r = \frac{n_1 \dots n_k}{m_1 \dots m_l}$$

n's, m's prime, not repeated.

Then  $p m_1^2 \dots m_l^2 = n_1^2 \dots n_k^2$   
 odd # of primes = even # of prime. — Lypus

∴ 'defect of beauty': decimal representations of real #'s are not unique:

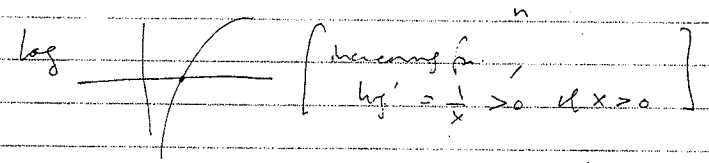
$$\boxed{.999\dots = 1.000\dots} \quad ?$$

$$\frac{9}{10} + \frac{9}{10^2} + \dots = \sum_{k \geq 1} 9 \cdot 10^{-k} = 9 \sum_{k \geq 0} 10^{-k-1}$$

$p \neq 2$   
 $4 \neq 2$   
 $8 \neq 2$   
 $23 \neq 7$

Claim  $\forall \epsilon < x < 1$  then  $x^n \rightarrow 0$

More precisely: For any  $\epsilon > 0$ ,  $\exists n$  such that  $x^n < \epsilon$



$$\epsilon > x^n \Leftrightarrow \log \epsilon > n \log x$$

$$- \log \epsilon < -n \log x$$

$$n > \frac{-\log \epsilon}{-\log x}$$

$$\boxed{n = \frac{-\log \epsilon}{-\log x}}$$

$\forall |x| < 1$   
 $h = n$   
 $\sum_{k=0}^{\infty} x^k \rightarrow \frac{1}{1-x} = \frac{x^{n+1}}{1-x} \rightarrow \frac{1}{1-x}$

ex  $x = \frac{1}{10}$

Lemma  $(1+x+\dots+x^n)(1-x) = 1-x^{n+1}$

$$\frac{1+x+\dots+x^n}{1-x} = \frac{1-x^{n+1}}{1-x}$$

"Telescoping sum"

$$= 1+x+\dots+x^n = \frac{1-x^{n+1}}{1-x} = \frac{1}{1-x} - \frac{1}{1-x} x^{n+1}$$

$$\lim_{n \rightarrow \infty} (a + b_n) = a + \lim_{n \rightarrow \infty} b_n = a$$

$$.999\dots = \frac{9}{10} \frac{1}{1-1/10} = 1.00\dots$$

? More elementary argument?

inf  $\{x^n \mid n=1,2,\dots\} = \text{greatest lower bound} = B$ ,

$$xB = x \cdot \inf \{x^n \mid n=1,2,\dots\} = \inf \{x^{n+1} \mid n=2,3,\dots\}$$

$$\geq \inf \{x^n \mid n=2,3,\dots\} = B$$

BUT  $B > xB \geq B$  if  $B \neq 0!$

11)

$$B = \inf \{x \mid x \in S\} \quad a > 0$$

$$\Rightarrow aB = \inf \{ax \mid x \in S\}$$

Proof

$$1) \quad x \in S \Rightarrow B \leq x$$

$$2) \quad (x \in S \Rightarrow B' \leq x) \Rightarrow (B' \leq B)$$

$a > 0$

$$1) \quad x \in S \Rightarrow C \leq ax \Rightarrow \vec{a}C \leq x$$

$$2) \quad (x \in S \Rightarrow C' \leq ax) \Rightarrow C' \leq C$$
  
$$\vec{a}C' \leq \vec{a}C$$

$$\Rightarrow \vec{a}C = B$$

Real numbers contain rational #'s,

$S = \{x \mid x^2 < 2\}$  below has least upper bound

$$\text{Ex } \sqrt{2} = \inf \{r \in \mathbb{Q} \mid r^2 < 2\}$$

Definition:  $r = \sup \{S \subset \mathbb{Q} \mid \text{if } s' \in \mathbb{Q}, s' > s \text{ then } s' \notin S\}$

x)

Blame Socrates! next initiate use

$\mathbb{N} = 0, 1, \dots$  = natural numbers

von Neumann:  $0 = \{\emptyset\}$  successor.

$$n+1 = \{n, \{n\}\} \quad 1 = \{0, \{0\}\}$$

Are in a kind of conceptual box, which means a way  
 $S \notin S, S \subset S$  we have stronger meaning  
of  $S$  in it.

Subsets are boxes containing collections  
of things in the box containing  $S$ .

$$\forall m+n \in \mathbb{N} \quad \exists m+n \in \mathbb{N}$$

$$l + (m+n) = (l+m) + n, \text{ etc.}$$

$$\exists \mathbb{Z} = \dots -3, -2, -1, 0, 1, 2, \dots$$

$$\forall k \in \mathbb{Z} \exists -k \Rightarrow k + (-k) = 0, \text{ etc.}$$

$\mathbb{Q}$  = rational numbers =

equivalence classes  $\mathbb{Z} \times (\mathbb{Z} - \{\emptyset\})$

$$(n, m) \rightarrow n/m$$

183 p54

$A, B$  sets  
 $(a, b)$   
 $A \times B = \{ \{a, \{a, b\}\} \mid a \in A, b \in B \}$

$$(a, b) = (c, d)$$

$(\Rightarrow)$

$$a, \{a, b\} \Leftrightarrow c, \{c, d\}$$
  
 $a \in \{a, b\}, c \in \{c, d\}$

$$\exists \text{ eq. relation } (n, m) \sim (a, b) \Leftrightarrow$$

$\exists j, k \in \mathbb{Z} \neq 0$  such that

$$kn = ja$$

$$km = jb$$

$$(6, 9) \sim (8, 12)$$

$\{a, b\} = c, \dots$   
trouble